TWINX

1* Introduction.   This short program compiles a master index for a set of programs that have been processed by CTWILL. To use it, you say, e.g., `twinx *.tex >index.tex`. The individual programs should define their names with a line of the form '`\def\title{NAME}`'.

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <ctype.h>
```
  ⟨ Type definitions 4* ⟩
  ⟨ Global variables 2 ⟩
  ⟨ Procedures 5* ⟩

  int $main(\textbf{int } argc, \textbf{char } *argv[\,])$
  {
    ⟨ Local variables 9* ⟩;
    ⟨ Initialize the data structures 8 ⟩;
    while $(--argc)$ {
      if $((f \leftarrow fopen(*{+}{+}argv, \texttt{"r"})) \equiv \Lambda)$
        $fprintf(stderr, \texttt{"twinx:}_{\sqcup}\texttt{Couldn't}_{\sqcup}\texttt{open}_{\sqcup}\texttt{file}_{\sqcup}\texttt{\%s}_{\sqcup}\texttt{for}_{\sqcup}\texttt{reading!\textbackslash n"}, *argv);$
      else {
        ⟨ Scan file f until coming to the title 3* ⟩;
        $fclose(f);\ memcpy(*argv + strlen(*argv) - 3, \texttt{"idx"}, 3);$
        if $((f \leftarrow fopen(*argv, \texttt{"r"})) \equiv \Lambda)$
          $fprintf(stderr, \texttt{"twinx:}_{\sqcup}\texttt{Couldn't}_{\sqcup}\texttt{open}_{\sqcup}\texttt{file}_{\sqcup}\texttt{\%s}_{\sqcup}\texttt{for}_{\sqcup}\texttt{reading!\textbackslash n"}, *argv);$
        else {
          ⟨ Copy the index file f into the data structures 10* ⟩;
          $fclose(f);$
        }
      }
    }
    ⟨ Output the data structures to make a master index 13 ⟩;
    return 0;
  }

3* For your convenience, TWINX grabs the first "word" in `\title` and turns it into uppercase form.

⟨ Scan file f until coming to the title 3* ⟩ ≡
  while (1) {
    if $(fgets(buf, buf\_size, f) \equiv \Lambda)$ {
      $fprintf(stderr, \texttt{"twinx:}_{\sqcup}\texttt{(no}_{\sqcup}\texttt{title}_{\sqcup}\texttt{found}_{\sqcup}\texttt{in}_{\sqcup}\texttt{file}_{\sqcup}\texttt{\%s)\textbackslash n"}, *argv);\ title[0] \leftarrow \texttt{'\textbackslash 0'};$ **break**;
    }
    if $(strncmp(buf, \texttt{"\textbackslash\textbackslash def\textbackslash\textbackslash title\textbackslash\{"}, 11) \equiv 0)$ { **char** $*p, *q;$
      for $(p \leftarrow buf + 11, q \leftarrow title;\ *p \wedge *p \neq \texttt{'}_{\sqcup}\texttt{'} \wedge *p \neq \texttt{'\}'};\ p{+}{+})\ *q{+}{+} \leftarrow toupper(*p);$
      $*q \leftarrow \texttt{'\textbackslash 0'};$ **break**;
    }
  }
This code is used in section 1*.

4*   Data structures.   Our main task is to collate a bunch of texts associated with keys that have already
been sorted. It seems easiest to do this by repeatedly merging the new data into the old, even though this
means we'll be passing over some of the same keys 30 times or more; the computer is fast, and this program
won't be run often.

Further examination shows that a merging strategy isn't so easy after all, because the sorting done by
CTWILL (and by CWEAVE) is weird in certain cases. When two index entries agree except for their "ilk," the
order in which they appear in the index depends on the order in which they appear in the program. Thus,
they might well appear in different order in two of the indexes we are merging. (There's also another glitch,
although not quite as devasting: When two index entries have the same letters and the same ilk, but differ
with respect to uppercase versus lowercase, the order in which they appear depends on the hash code used
in CWEB's common.w code!)

So we'll use Plan B: All index entries will first be copied into a long list. The list will almost always consist
of many sorted sublists, but we will not assume anything about its order. After all the copying has been
done, we will use a list-merge sort to finish the job.

The data structure is built from nodes that each contain three pointers. The first pointer is to an *id* string;
the third pointer is to the *next* node; and the second pointer is either *data.s*, a pointer to a string of text,
or *data.n*, a pointer to a node. In the main list, the *id* fields are the keys of the index, and the *data.n* fields
point to lists of associated texts. In the latter lists, the *id* fields are the individual program titles, while the
data.s fields are the texts.

⟨ Type definitions 4* ⟩ ≡

  typedef union {
    char ∗*s*;
    struct node_struct ∗*n*;
  } **mixed**;
  typedef struct node_struct {
    const char ∗*id*;
    mixed *data*;
    struct node_struct ∗*next*;
  } **node**;

This code is used in section 1*.

5* We copy strings into blocks of storage that are allocated as needed. Here's a routine that stashes away a given string. It makes no attempt to handle extremely long strings, because such strings will arise only if the input is all screwed up.

#**define** *string_block_size* 8192   ▷ number of bytes per string block ◁

⟨ Procedures 5* ⟩ ≡

  char **save_string*(**char** **s*)
  {
    char **p, *q*;
    int *l*;
    for (*p ← s*; **p*; *p++*) ;
    l ← p − s + 1;
    if (*l > string_block_size*) {
      fprintf(*stderr*, "twinx:␣Huge␣string␣'%.20s...'␣will␣be␣truncated!\n", *s*);
      l ← string_block_size; *s[l − 1] ← '\0'*;
    }
    if (*next_string + l ≥ bad_string*) {
      *next_string ← (**char** *) malloc*(*string_block_size*);
      if (*next_string ≡ Λ*) {
        fprintf(*stderr*, "twinx:␣Not␣enough␣room␣for␣strings!\n"); *exit*(−1);
      }
      bad_string ← next_string + string_block_size;
    }
    for (*p ← s, q ← next_string*; **p*; *p++*) **q++ ← *p*;
    *q ← '\0'; *next_string ← q + 1*; **return** *next_string − l*;
  }

See also sections 6*, 17*, and 20*.

This code is used in section 1*.

6* Nodes are allocated with a similar but simpler mechanism.

#**define** *nodes_per_block* 340

⟨ Procedures 5* ⟩ +≡

  node **new_node*(**void**)
  {
    if (*next_node ≡ bad_node*) {
      *next_node ← (**node** *) calloc*(*nodes_per_block*, **sizeof**(**node**));
      if (*next_node ≡ Λ*) {
        fprintf(*stderr*, "twinx:␣Not␣enough␣room␣for␣nodes!\n"); *exit*(−2);
      }
      bad_node ← next_node + nodes_per_block;
    }
    next_node ++; **return** *next_node − 1*;
  }

9* ⟨ Local variables 9* ⟩ ≡

  node **main_node*;   ▷ current end of main list ◁

This code is used in section 1*.

10* Copying. Lines in the index file f that we're reading either begin a new entry or continue a long entry. In the first case, the line begins with \I and then either \\{key} or \|{key} or \.{key} or \&{key} or \${key} or \9{key} or just ␣␣{key}. (These correspond to multi-character italic, single-digit italic, typewriter, bold, custom, variable, and roman styles.) In the second case, the line begins with a page number or \[; however, we recognize the second case by the fact that the previous line did not end with a period.

⟨ Copy the index file f into the data structures 10* ⟩ ≡
  while (1) { **node** *cur_node;

    if ($fgets(buf, buf\_size, f) \equiv \Lambda$) **break**;   ▷ end of file ◁
    if ($strncmp(buf, "\\I", 2) \equiv 0$) {
      ⟨ Copy a new index entry into cur_name and cur_node 11* ⟩;
      $main\_node{\rightarrow}next \leftarrow new\_node(\,)$; $main\_node \leftarrow main\_node{\rightarrow}next$;
      $main\_node{\rightarrow}id \leftarrow save\_string(cur\_name)$; $main\_node{\rightarrow}data.n \leftarrow cur\_node$;
    }
    else if ($buf[0] \neq$ '\n')
      $fprintf(stderr, $"twinx:␣couldn't␣deal␣with␣'%.10s...'␣in␣file␣%s!\n"$, buf, {*}argv)$;
  }

This code is used in section 1*.

11* ⟨ Copy a new index entry into cur_name and cur_node 11* ⟩ ≡
  if ($buf[4] \neq$ '{') {
    $fprintf(stderr, $"twinx:␣missing␣brace␣in␣file␣%s:␣'%.20s...'\n"$, {*}argv, buf)$; **break**;
  }
  { **char** *p, *q; **int** $bal \leftarrow 1$;

    $cur\_name[0] \leftarrow buf[2]$; $cur\_name[1] \leftarrow buf[3]$; $cur\_name[2] \leftarrow$ '{';
    for ($p \leftarrow buf + 5, q \leftarrow cur\_name + 3$; $*p \wedge (bal \vee *p \equiv$ '{'$)$; $p{+}{+}$) {
      switch ($*p$) {
      case '\\': $*q{+}{+} \leftarrow *p{+}{+}$; **break**;
      case '{': $bal{+}{+}$; **break**;
      case '}': $bal{-}{-}$; **break**;
      }
      $*q{+}{+} \leftarrow *p$;
    }
    if (bal) {
      $fprintf(stderr, $"twinx:␣unbalanced␣entry␣in␣file␣%s:␣'%.20s...'\n"$, {*}argv, buf)$; **break**;
    }
    if ($*p{+}{+} \neq$ ',') {
      $fprintf(stderr, $"twinx:␣missing␣comma␣in␣file␣%s:␣'%.20s...'\n"$, {*}argv, buf)$; **break**;
    }
    if ($*p{+}{+} \neq$ '␣') {
      $fprintf(stderr, $"twinx:␣missing␣space␣in␣file␣%s:␣'%.20s...'\n"$, {*}argv, buf)$; **break**;
    }
    $*q \leftarrow$ '\0'; ⟨ Copy the text part of the index entry into cur_node 12* ⟩;
  }

This code is used in section 10*.

12* When we get here, p points to the beginning of the text following a key in the index. The index entry ends with the next period, possibly several lines hence. In the multiple-line case, *cur_node* will point to the final line, which points to the penultimate line, etc.

⟨ Copy the text part of the index entry into *cur_node* 12* ⟩ ≡
```
{ int period_sensed ← 0;
  node *continuation;

  cur_node ← new_node( );  cur_node→id ← save_string(title);
  do {
    for (q ← p; *q ∧ *q ≠ '\n' ∧ *q ≠ '.'; q++) ;
    if (*q ≡ '.') period_sensed ← 1;
    *q ← '\0';  cur_node→data.s ← save_string(p);
    if (period_sensed) break;
    continuation ← new_node( );       ▷ the id field is Λ ◁
    continuation→next ← cur_node;  cur_node ← continuation;  p ← buf;
  } while (fgets(buf, buf_size, f));
  if (¬period_sensed) {
    fprintf(stderr, "twinx:␣File␣%s␣ended␣in␣middle␣of␣entry␣for␣%s!\n", *argv, cur_name);
    break;
  }
}
```
This code is used in section 11*.

14* The *compare* subroutine, which specifies the relative order of *id* fields in two nodes, appears below. Let's get the sorting logic right first.

The algorithm is, in fact, rather pretty—I hate to say cute, but that's the word that comes to mind. Some day I must write out the nice invariant relations in these loops. Too bad it's not more efficient.

Remember that *header.id* is $-\infty$ and *sentinel.id* is $+\infty$. Also remember that the main list begins and ends at the header node.

⟨ Sort the main list, collapsing entries with the same *id* 14* ⟩ ≡
```
main_node→next ← &header;
while (1) { node *p, *q, *r, *s, *t;

  t ← &header;  r ← t→next;
  while (1) {
    if (r ≡ &header) break;
    p ← s ← r;  ⟨ Advance s until it exceeds r ← s→next 15* ⟩;
    if (r ≡ &header) break;
    s→next ← &sentinel;  q ← s ← r;  ⟨ Advance s until it exceeds r ← s→next 15* ⟩;
    s→next ← &sentinel;  ⟨ Merge p and q, appending to t 16* ⟩;
    t→next ← r;
  }
  if (t ≡ &header) break;
}
```
This code is used in section 13.

15* ⟨ Advance s until it exceeds $r \leftarrow s \rightarrow next$ 15* ⟩ ≡
  do { **int** d;

    $r \leftarrow s \rightarrow next$; $d \leftarrow compare(s, r)$;
    if $(d > 0)$ **break**;   ▷ $s \rightarrow id > r \rightarrow id$ ◁
    if $(d \equiv 0)$ {   ▷ $s \rightarrow id \leftarrow r \rightarrow id$ ◁
      $collapse(s, r)$;   ▷ put r's data into s's list ◁
      $s \rightarrow next \leftarrow r \rightarrow next$;   ▷ node r will be unclaimed garbage ◁
    }
    else $s \leftarrow r$;   ▷ this is the normal case, $s \rightarrow id < r \rightarrow id$ ◁
  } **while** (1);

This code is used in section 14*.

16* Merging takes place in such a way that sorting is stable. Thus, index entries for a key that appears in different programs will remain in the order of the `.tex` files on the command line.

⟨ Merge p and q, appending to t 16* ⟩ ≡
  do { **int** d;

    $d \leftarrow compare(p, q)$;
    if $(d > 0)$ {   ▷ $p \rightarrow id > q \rightarrow id$ ◁
      $t \rightarrow next \leftarrow q$; $t \leftarrow q$; $q \leftarrow q \rightarrow next$;
    }
    else if $(d < 0)$ {   ▷ $p \rightarrow id < q \rightarrow id$ ◁
      $t \rightarrow next \leftarrow p$;   ▷ $p \rightarrow id < q \rightarrow id$ ◁
      $t \leftarrow p$; $p \leftarrow p \rightarrow next$;
    }
    else if $(p \equiv \& sentinel)$ **break**;
    else {
      $collapse(p, q)$;   ▷ put q's data into p's list ◁
      $q \leftarrow q \rightarrow next$;
    }
  } **while** (1);

This code is used in section 14*.

17* Comparison is a three-stage process in general. First we compare the keys without regarding case or format type. If they are equal with respect to that criterion, we try again, with case significant. If they are still equal, we look at the format characters (the first two characters of the id field).

⟨ Procedures 5* ⟩ +≡
  int $compare(\mathbf{node} *p, \mathbf{node} *q)$
  { **unsigned char** $*pp, *qq$;

    for $(pp \leftarrow (\mathbf{unsigned\ char} *) p \rightarrow id + 3, qq \leftarrow (\mathbf{unsigned\ char} *) q \rightarrow id + 3;\ *pp \wedge ord[*pp] \equiv ord[*qq];$
        $pp ++, qq ++)$;
    if $(*pp \vee *qq)$ **return** $ord[*pp] - ord[*qq]$;
    for $(pp \leftarrow (\mathbf{unsigned\ char} *) p \rightarrow id + 3, qq \leftarrow (\mathbf{unsigned\ char} *) q \rightarrow id + 3;\ *pp \wedge *pp \equiv *qq;$
        $pp ++, qq ++)$;
    if $(*pp \vee *qq)$ **return** $(\mathbf{int}) *pp - (\mathbf{int}) *qq$;
    if $(p \rightarrow id[0] \neq q \rightarrow id[0])$ **return** $p \rightarrow id[0] - q \rightarrow id[0]$;
    return $p \rightarrow id[1] - q \rightarrow id[1]$;
  }

19* The right brace is placed lowest in collating order, because each key is actually followed by a right brace when we are sorting.

Apology: I haven't had time to update this part of the program to allow 8-bit characters. At present the data is assumed to be 7-bit ASCII, as it was in the early versions of CWEAVE.

⟨Initialize the data structures 8⟩ +≡
  $collate[0] \leftarrow 0$; $strcpy(collate + 1,$
      `"}`␣`\1\2\3\4\5\6\7\10\11\12\13\14\15\16\17\20\21\22\23\24\25\26\27\30\31\32\33\34\`
      `\35\36\37!\42#$%&'()*+,-./:;<=>?@[\\]^'{|~_abcdefghijklmnopqrstuvwxyz0123456789")`;
  `{ int` j;
    for $(j \leftarrow 1;\ collate[j];\ j{+}{+})\ ord[(\textbf{int})\,collate[j]] \leftarrow j$;
    $ord[128] \leftarrow j$;   ▷ this affects the ordering of $sentinel.id$ ◁
    for $(j \leftarrow \text{'A'};\ j \leq \text{'Z'};\ j{+}{+})\ ord[j] \leftarrow ord[tolower(j)]$;
  `}`

20* When two lists are combined, we put the data from the second node before the data from the first node, because we are going to reverse the order when printing. After this procedure has acted, the field $q{\rightarrow}data.n$ should not be considered an active pointer.

⟨Procedures 5*⟩ +≡
  void $collapse(\textbf{node}\ {*}p, \textbf{node}\ {*}q)$
  `{ node` ${*}x$;
    for $(x \leftarrow q{\rightarrow}data.n;\ x{\rightarrow}next;\ x \leftarrow x{\rightarrow}next)$;
    $x{\rightarrow}next \leftarrow p{\rightarrow}data.n$; $p{\rightarrow}data.n \leftarrow q{\rightarrow}data.n$;
  `}`

21* The only remaining trick is to format the underline characters properly, especially in the "custom" format when they must become x's.

⟨Output the main list in suitable TEX format 21*⟩ ≡
  `{ node` ${*}x$;
  $printf(\texttt{"\\\\input}␣\texttt{cttwinxmac\\n"})$;
  for $(x \leftarrow header.next;\ x \neq \&header;\ x \leftarrow x{\rightarrow}next)$ {
    $printf(\texttt{"\\\\I"})$; ⟨Output $x{\rightarrow}id$ in suitable TEX format 22*⟩;
    ⟨Output the lines of $x{\rightarrow}data.n$ in reverse order 23*⟩;
  `}`
  $printf(\texttt{"\\\\fin\\n"})$;
  `}`

This code is used in section 13.

22* ⟨Output $x{\rightarrow}id$ in suitable TEX format 22*⟩ ≡

```
{ const char *p ← x→id;
    if (*p ≡ '␣') {
        if (*(p + 1) ≠ '␣') goto unknown;
        goto known;
    }
    if (*p ≠ '\\') goto unknown;
    switch (*(p + 1)) {
    case '\\': case '|': case '.': case '&': case '9': printf("\\%c", *(p + 1)); goto known;
    case '$': printf("$\\");
        for (p += 3; *p ≠ '}'; p++)
            if (*p ≡ '_') putchar('x');
            else putchar(*p);
        putchar('$'); goto done;
    default: goto unknown;
    }
unknown: fprintf(stderr, "twinx:␣'%s'␣has␣unknown␣format!\n", p);
known:
    for (p += 2; *p; p++) {
        if (*p ≡ '_' ∧ *(p − 1) ≠ '\\') putchar('\\');
        putchar(*p);
    }
done: ;
}
```

This code is used in section 21*.

23* ⟨Output the lines of $x{\rightarrow}data.n$ in reverse order 23*⟩ ≡

```
{ node *y ← x→data.n, *z ← Λ;
    while (y) { node *w;
        w ← y→next; y→next ← z; z ← y; y ← w;
    }
    while (z) {
        if (z→id) printf("\\unskip,␣{\\sc␣%s}~", z→id);
        fputs(z→data.s, stdout); z ← z→next;
        if (z) putchar('\n');
        else puts(".");
    }
}
```

This code is used in section 21*.

24.* Index.

The following sections were changed by the change file: 1, 3, 4, 5, 6, 9, 10, 11, 12, 14, 15, 16, 17, 19, 20, 21, 22, 23, 24.

⟨ Advance *s* until it exceeds *r* ← *s⇁next* 15* ⟩   Used in section 14*.
⟨ Copy a new index entry into *cur_name* and *cur_node* 11* ⟩   Used in section 10*.
⟨ Copy the index file *f* into the data structures 10* ⟩   Used in section 1*.
⟨ Copy the text part of the index entry into *cur_node* 12* ⟩   Used in section 11*.
⟨ Global variables 2, 7, 18 ⟩   Used in section 1*.
⟨ Initialize the data structures 8, 19* ⟩   Used in section 1*.
⟨ Local variables 9* ⟩   Used in section 1*.
⟨ Merge *p* and *q*, appending to *t* 16* ⟩   Used in section 14*.
⟨ Output the data structures to make a master index 13 ⟩   Used in section 1*.
⟨ Output the lines of *x⇁data.n* in reverse order 23* ⟩   Used in section 21*.
⟨ Output the main list in suitable TEX format 21* ⟩   Used in section 13.
⟨ Output *x⇁id* in suitable TEX format 22* ⟩   Used in section 21*.
⟨ Procedures 5*, 6*, 17*, 20* ⟩   Used in section 1*.
⟨ Scan file *f* until coming to the title 3* ⟩   Used in section 1*.
⟨ Sort the main list, collapsing entries with the same *id* 14* ⟩   Used in section 13.
⟨ Type definitions 4* ⟩   Used in section 1*.